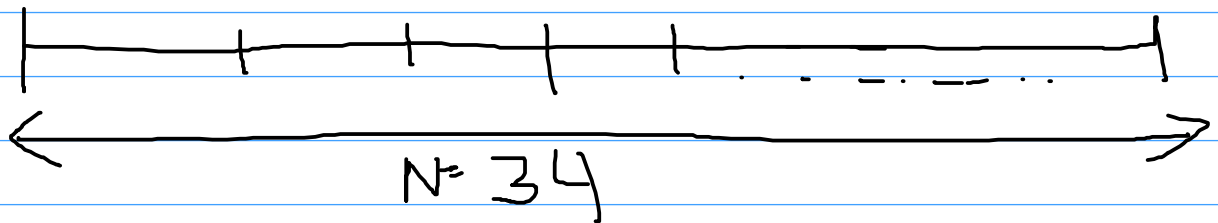


## Meet in the Middle

For a large sequence of data, some processing like calculating the sum of subsets is quite expensive in computation.

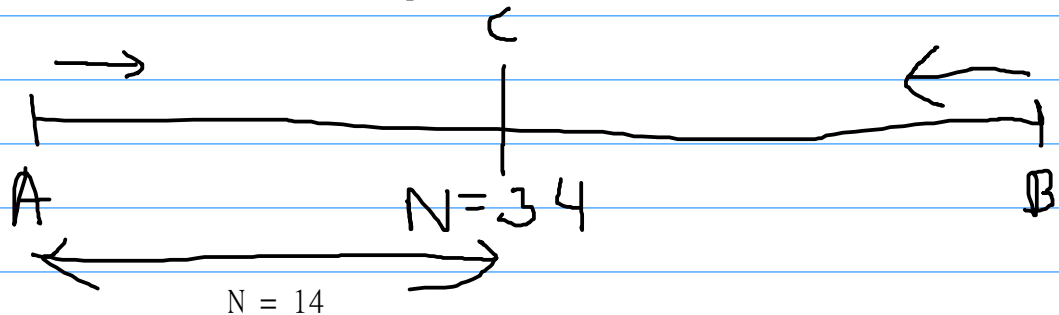


if we calculate the all the possible sum of subset for the above example, then

$$\text{No. of subset} = 2^{34} \approx 10^{11}$$

1 sec.  $\rightarrow 10^7$  computation.

After Meet in the Middle concept:

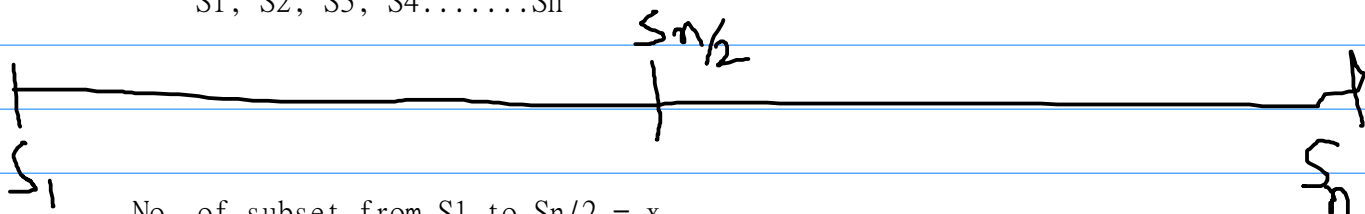


$$\text{No. of subset} = 2^{14}$$

$$\text{No. of computation} = 2^{14} + 2^{14} = 2^{15} \approx 10^5$$

Question Link : <https://www.spoj.com/problems/SUBSUMS/>

$S_1, S_2, S_3, S_4, \dots, S_n$



No. of subset from  $S_1$  to  $S_{n/2} = x$

No. of subset from  $S_{(n/2+1)}$  to  $S_n = y$

Lets take one subset sum from  $S_1$  to  $S_{n/2}$  region and same for  $S_{(n/2+1)}$  to  $S_n$  region.

$a+b$  = Sum of the subset from whole set i.e.  $S_1$  to  $S_n$ .

$$A \leq a+b \leq B$$

$$A - a \leq b \leq B - a$$

$$b \in [A - a, B - a + 1)$$

Lower bound

Upper bound

.